Resource Management for Next-generation HPC Systems: Challenges and Solutions

11th Scheduling for Large-scale Systems Workshop

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LLNL HPC Systems

| System | Processor | Nodes Cores | | Peak | |
|-----------------------|--------------------------|-------------|--------|-----------|-----|
| (Program) | Architecture | | | (TFLOP/s) | |
| RZ | | | | | |
| RZCereal (M&IC) | Intel Xeon E5530 | 21 | 169 | 1.6 | |
| <u>RZHasGPU</u> | Intel Xeon E5-2667 v3 | 20 | 320 | 8.2 | |
| RZMerl (ASC/ M&IC) | Intel Xeon E5-2670 | 162 | 2,592 | 53.9 | |
| RZSLIC *** | Intel Xeon E5330 | 3 | Stats | | |
| RZuSeq (ASC) **** | IBM PowerPC A2 | 522 | Max: 9 | 8,304 no | des |

267

Intel 2

X566 **AMD**

8354 Intel >

E5-2670

X5660

A2

Intel Xeon EP

IBM PowerPC

Intel Xeon

E5-2670

2,916

46,656

| tems | | | System (Program) | Processor Architecture | Nodes | Cores | Peak (TFLOP/s) | |
|---|--------|-----------|------------------------|---------------------------|--------|--------|-------------------|---|
| | | | CZ | | | | | |
| ores Peak (TFLOP/s) | |) | Ansel (M&IC) | Intel Xeon EP X5660 | 324 | 3,888 | 43.5 | ŀ |
| | | | Aztec (M&IC) | Intel Xeon EP X5660 | 96 | 1,152 | 12.9 | |
| 169 | 1.6 | | Catalyst (ASC/M&IC) | Intel Xeon E5-2695 v2 | 324 | 7,776 | 149.3 | |
| 320 | 8.2 | | Cab (ASC/ M&IC) | Intel Xeon E5-2670 | 1,296 | 20,736 | 431.3 | |
| 2,592 | 53.9 | | | MD Ontoron | | | | |
| State | 5 | | | | | 256 | 1.6 | |
| Max: 98,304 nodes in one system (Sequoia) | | | | | | | | |
| 25 systems across open and closed zones | | | | | | 112.7 | | |
| · | | | | | 40 | _ | | |
| Various processor architectures | | | | | 23,328 | 261.3 | | |
| 670 | 50Z | 7,007 | 107 | ntel Xeon | 162 | 2,592 | 53.9 | |
| Keon EP | 1,296 | 15,552 | 174.2 | :5-2670 | | | | |
| PowerPC | 98,304 | 1,572,864 | 20,132 | ntel Xeon 5-2670 | 324 | 5,056 | 107.8 | |

24,576

3M PowerPC

2

970.4



Intel Xeon

Inca (ASC)

Juno (ASC)

Max (ASC)

Muir (ASC)

<u>Sequoia</u>

(ASC) **

Zin (ASC)

E5530

RZZeus

(M&IC)



5,033

393,216

Future systems present several new challenges

- Run-to-run <u>variability</u>, inter-job interference
- Multiple constraints: power, network, I/O, and data awareness
- Need to support high-throughput workloads, such as UQ workloads
- Increased error and failure rates



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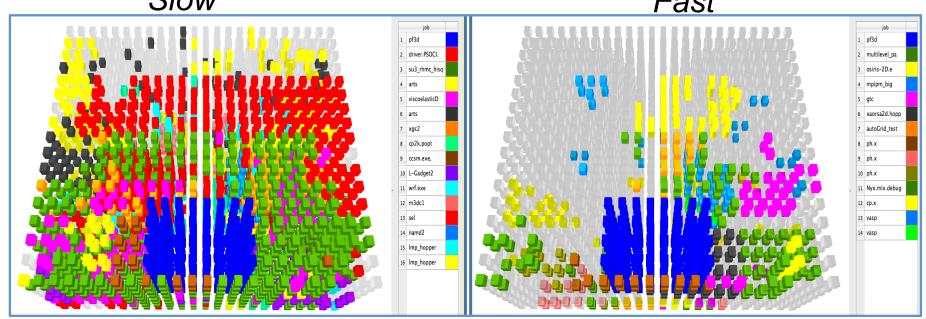
How do we design low-overhead, scalable resource managers?



PROBLEM: Network contention and inter-job interference can lead to severe run-to-run variability

Two 512-node pF3D runs (blue) on Hopper Cray XE6

Slow
Fast



25% slower messaging rate due to other jobs!

Graph Courtesy: Abhinav Bhatele, LLNL



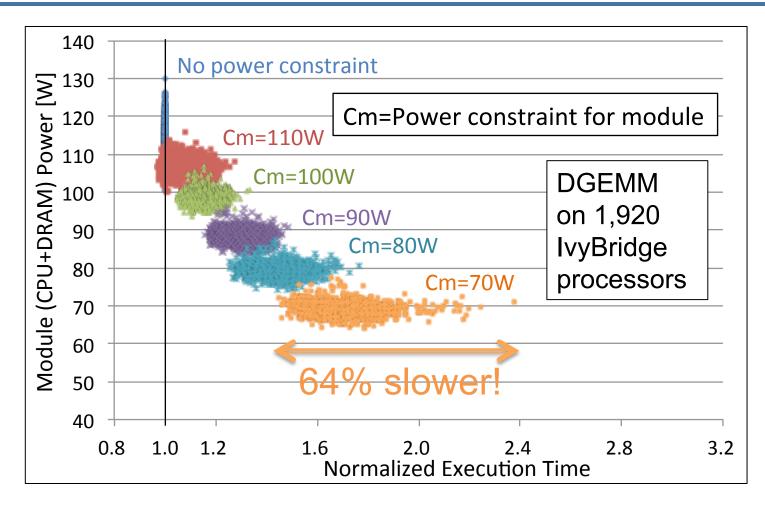


PROBLEM: Existing resource managers may not scale well for high-throughput ensemble workloads

- Managing several thousand jobs at once can be slow, e.g., UQ ensemble workloads
 - Job launch times vary
- Typically targeted toward fewer large-sized jobs, but not for more, small-sized jobs
- ~250 jobs on clusters that have thousands of nodes



PROBLEM: Power capping affects each processor differently, creating runtime variability and imbalance

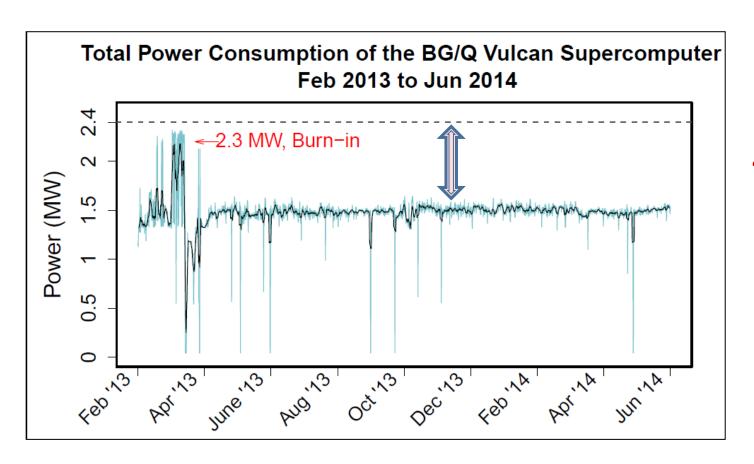


Inadomi et al., SC'15

RESULT: Unhappy Users!



UNDERUTILIZED CAPACITY: Utilizing one resource well leads to under-utilization of another resource



40%
Procured
Power
Unused!

Patki et al., HPDC'15

Existing resource managers cannot be extended to support multi-constraint HPC systems easily

| | Network Topology | Network Bandwidth | I/O- Aware | Power- Aware | Type? |
|--------------------|---------------------|----------------------|---------------|-----------------|-------------------------|
| SLURM | ✓ | X | X | • | Monolithic |
| Moab/ Torque | • | × | X | × | Monolithic |
| LSF, IBM | ✓ | X | • | • | Monolithic |
| Cobalt, Argonne | ✓ | × | X | × | Monolithic |
| Mesos, Apache | ~ | × | X | • | 2-Level Hierarchical |
| PBSPro, Altair | • | × | × | × | Monolithic |

: Limited support,In progress



Existing resource managers are not designed to be fault-tolerant

- Moldable and malleable jobs are typically not supported
- Checkpoint/Restart process is slow and unintelligent
- Users end up requesting 'redundant' nodes and more time as part of allocation
 - Underutilized capacity (see Felix's talk!)

RESULT: Unhappy System Administrators!



LLNL's approach is to provide a holistic solution for a large-scale HPC system

Monitoring infrastructure for production clusters

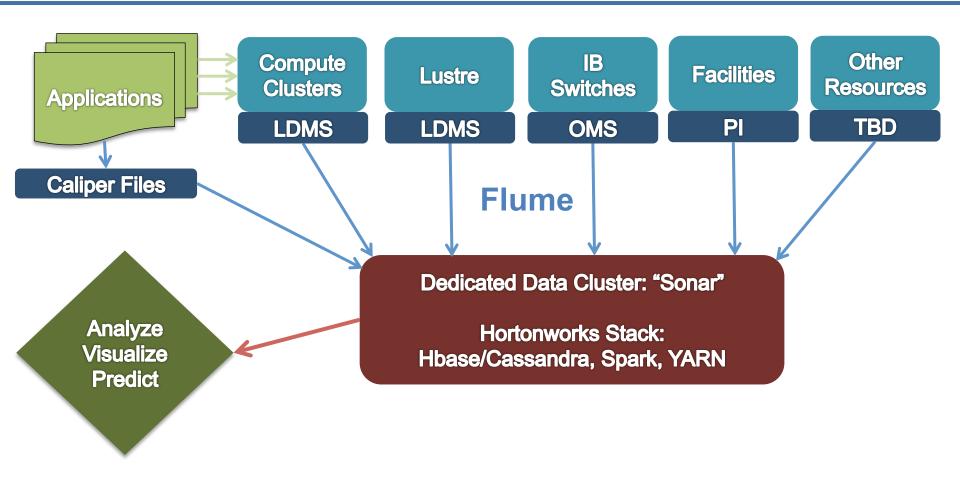
PIPER*, SONAR cluster

Framework for nextgen resource management



*Performance Insights for Programmers and Exascale Runtimes

Deploying site-wide monitoring infrastructure



PI: Todd Gamblin, LLNL

Flux Framework: Next Generation Resource Mgmt.

- Hierarchical resource manager designed to support future HPC systems in a scalable manner
- Three key components: flux-core, flux-sched, flux-capacitor
- Open source, initial release will be available soon

Flux Framework: flux-core

- Communication layer comprising of distributed message broker and plug-in modules for services
- Three overlay networks implemented using ZeroMQ
 - For request/response, session-wide broadcasts, and debugging
- Flux KVS/DHT module for job and resource configuration in a session
 - Useful for logging, synchronization, broadcasting

Flux Framework: flux-sched

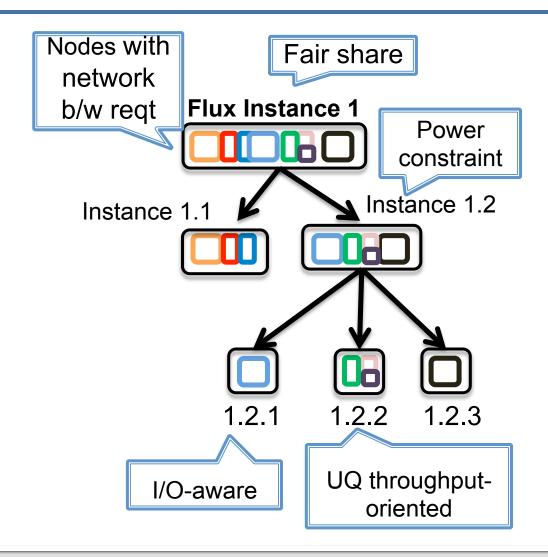
- Flexible resource model to represent new hardware and flow resources
 - Accelerators
 - Power, network bandwidth
 - Easily extensible and mutable
- Scalable job management with asynchronous eventbased protocols
- Currently supports FCFS, easy backfilling



Flux Framework: flux-capacitor

- Simple python interface for high-throughput, ensemble workloads
- Supports pulsed job launch: feeds jobs (flux instances) to the system at an ideal rate

Flux Framework: Example



Flux Framework: Current Research

- Hierarchical Scheduling
 - Adds additional levels of schedulers to form a hierarchy
 - Initial study shows adding only one additional level in the reduces the scheduling complexity by 3.6x
 - But increases the resource fragmentation by up to 20%
- Dynamic Scheduling
 - Allows for allocations to change size at runtime
 - Fault-tolerance, fragmentation, utilization
- I/O Aware Scheduling
 - Efficient ways to schedule for systems with burst buffers
 - Increases system efficiency by ~1.3x in exchange for increasing turnaround time by ~1.5x



Open Research Questions

- Sources of run-to-run variability
- Analysis of user behavior, workloads, error logs
- Impact of data staging on power and network performance
- How to prioritize constraints: power, network, file systems

The Flux Team





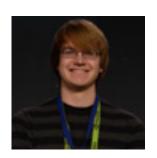








Lawrence Livermore National Laboratory*



University of Delaware





*And Mark Grondona, whose picture I couldn't find this morning!

Flux Resources



Project Page: <u>flux-framework.github.io</u>

flux-core: github.com/flux-framework/flux-core

flux-sched: github.com/flux-framework/flux-sched

flux-capacitor: github.com/flux-framework/capacitor



